**Conflict Serializability**

A schedule is conflict serializable if it is conflict equivalent to a serial schedule.

|  |  |
| --- | --- |
| **Non-conflict pairs:**  R(A) R(A)  R(B) R(A)  W(B) R(A)  R(B) W(A)  W(A) W (B) | **Conflict pairs:**  R(A) W(A)  W(A) R (A)  W(A) W (A) |

|  |  |
| --- | --- |
| T1 | T2 |
| R(A)  W(A)  R(B)  W(B) | R(A)  W(A)  R(B)  W(B) |

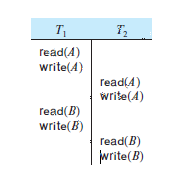
|  |  |
| --- | --- |
| T1 | T2 |
| R(A)  W(A)  R(B)  W(B) | R(A)  W(A)  R(B)  W(B) |

|  |  |
| --- | --- |
| T1 | T2 |
| R(A)  W(A)  R(B)  W(B) | R(A)  W(A)  R(B)  W(B) |

So, this schedule is serializable.

|  |  |
| --- | --- |
| T1 | T2 |
| Read (A)  Write(A)  R(B)  W(B) | R(A)  W(A)  R(B)  W(B) |

|  |  |
| --- | --- |
| T1 | T2 |
| Read (A)  Write(A)  R(B)  W(B) | R(A)  W(A)  R(B)  W(B) |



This is also conflict serializable.

So, 1st schedule is serializable.